

BISM: Bytecode-Level Instrumentation for Software Monitoring

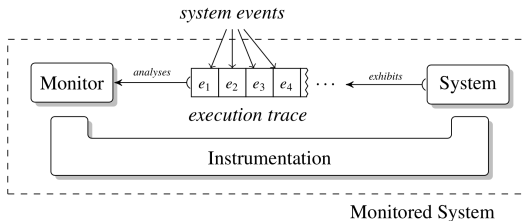
Chukri Soueidi Ali Kassem Yliès Falcone

Univ. Grenoble Alpes, Inria

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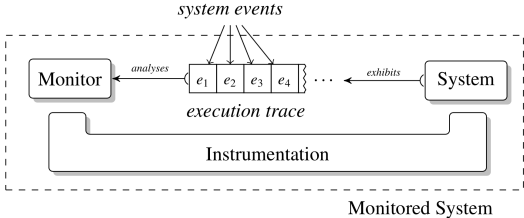


Instrumentation is essential in Software Monitoring



- Allows us to *extract* and *abstract* system events

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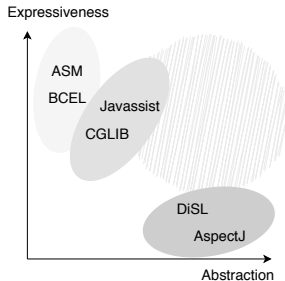


- Allows us to *extract* and *abstract* system events
- *Granularity level* of extracted events may range from:
 - Coarse (e.g., a method call)
 - Fine-grained (e.g., a jump in the control flow)

Instrumentation Tools

Several general-purpose tools for instrumenting Java programs varying in:

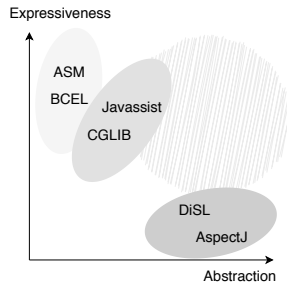
- Levels of *expressiveness*
- Levels of *abstraction*



Instrumentation Tools

Several general-purpose tools for instrumenting Java programs varying in:

- Levels of *expressiveness*
- Levels of *abstraction*



	BCEL	ASM	Javassist	CGLIB	DiSL	AspectJ
Bytecode Visibility	✓	✓	✓	✓	✓	✗
Allows bytecode insertion	✓	✓	✓	✓	✗	✗
No bytecode proficiency	✗	✗	✗	✗	✓	✓
High level of abstraction	✗	✗	✗	✗	✓	✓
Follows AOP paradigm	✗	✗	✗	✗	✓	✓

Motivating Example

- Extract fine-grained `eT(..)` and `eF(..)` events after conditional jumps
- Can we automate instrumentation?

```
g_authenticated = BOOL_FALSE;
if(g_ptc > 0) {

    int cmp = byteArrayCompare(...);
    if(cmp == BOOL_TRUE) {

        g_ptc = 3;
        g_authenticated = BOOL_TRUE;
    }
    else {

        g_ptc--;
    }
}
```

Original code

```
g_authenticated = BOOL_FALSE;
if(g_ptc > 0) {
    eT(1, ">", g_ptc, 0);
    int cmp = byteArrayCompare(...);
    if(cmp == BOOL_TRUE) {
        eT(2, "==", cmp, BOOL_TRUE);
        g_ptc = 3;
        g_authenticated = BOOL_TRUE;
    }
    else {
        eF(2, "==", cmp, BOOL_TRUE);
        g_ptc--;
    }
} else
    eF(1, ">", g_ptc, 0);
```

Instrumented code

Motivating Example

- Extract fine-grained `eT(..)` and `eF(..)` events after conditional jumps
- Can we automate instrumentation?

Yes, possible in DiSL but not with AspectJ

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g_authenticated = BOOL_FALSE;
if(g_ptc > 0) {

    int cmp = byteArrayCompare(...);
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Original code

```

g_authenticated = BOOL_FALSE;
if(g_ptc > 0) {
    eT(1, ">", g_ptc, 0);
    int cmp = byteArrayCompare(...);
    if(cmp == BOOL_TRUE) {
        eT(2, "==", cmp, BOOL_TRUE);
        g_ptc = 3;
        g_authenticated = BOOL_TRUE;
    }
    else {
        eF(2, "==", cmp, BOOL_TRUE);
        g_ptc--;
    }
} else
    eF(1, ">", g_ptc, 0);

```

Instrumented code

Another Motivating Example

- Duplicate all if statements and report inconsistency
- Can we automate instrumentation?

```
// a simple if statement
if(g_ptc > 0) {

    int cmp = compare(...);

}
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Original code

```
// duplicating the if statement
if(g_ptc > 0) {
    if(g_ptc > 0){
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    }
}
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    if(g_ptc > 0 ) {
        Report Attack!
    }
}
```

Instrumented code

Another Motivating Example

- Duplicate all if statements and report inconsistency
- Can we automate instrumentation?

Yes, but neither with AspectJ nor DiSL¹

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Instrumented code

¹A workaround can have the same effect (but with a lot more instrumented code) as we will see in the experiments.

Contribution: The BISM tool

BISM is a **lightweight** Java **bytecode** instrumentation tool that features an **expressive** and **high-level** instrumentation language

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- Features a **control-flow-aware** instrumentation language

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- Built-in **Transformers** (instrumentation classes)

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BISM is a **lightweight** Java **bytecode** instrumentation tool that features an **expressive** and **high-level** instrumentation language

- Follows the **aspect-oriented programming** (AOP) paradigm
- Features a **control-flow-aware** instrumentation language
- Built-in **Transformers** (instrumentation classes)
- Better **performance** than DiSL and AspectJ demonstrated in:
 - Two security scenarios with
 - Advanced Encryption Standard
 - Financial Transactions System
 - Classic RV scenario with DaCapo benchmarks

Outline

- 1 Motivation
- 2 BISM Design Goals & Features**
- 3 BISM Language
- 4 BISM Implementation
- 5 Evaluation
- 6 Conclusion

Design Goals & Features

■ Convenient instrumentation mechanism

- *Modularity* in separate instrumentation classes that follow AOP paradigm
- Granularity in joinpoints ranging from bytecode instruction to methods

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- Dynamic context objects (local variables, stack values, method args ...)
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- Out-of-the-box *CFG information* in static context objects
- *Conditional jumps* joinpoints (OnTrueBranch, OnFalseBranch)
- *Visualizer* for CFGs

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■ Compatibility with ASM

- Bytecode *instruction insertion* using ASM syntax
- Preserve ASM object structure

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■ Two Instrumentation modes

- *Build-time* for static instrumentation
- *Load-time* as a Java agent

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Joinpoints

Identify specific bytecode locations in target program and capture their execution

■ Instruction Joinpoints

- BeforeInstruction: *before* bytecode instruction
- AfterInstruction: *after* bytecode instruction
- BeforeMethodCall: *after loading* all needed values on stack
- AfterMethodCall: *before storing* return value

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■ Basic Block Joinpoints

- OnBasicBlockEnter: at *entry* of basic block
- OnBasicBlockExit: *after last instruction* of basic block
- OnTrueBranchEnter: *after conditional jump* on True evaluation
- OnFalseBranchEnter: *after conditional jump* on False evaluation

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■ Method Joinpoints

- OnMethodEnter: on *method entry* block
- OnMethodExit: on all *exit blocks* of method before return or throw

Static Context

Provide relevant static information at joinpoints

■ Instruction

- `index, opcode, next, previous, isConditionalJump()...`
- `node, getBasicValueFrame(), getSourceValueFrame()`
- `methodName, basicBlock, className`

■ MethodCall

- `methodOwner, methodName, currentClassName`
- `node, instruction`

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Provide relevant static information at joinpoints

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■ MethodCall

- `methodOwner, methodName, currentClassName`
- `node, instruction`

■ BasicBlock

- `id, index, blockType, size, getFirstInstruction() ...`
- `getSuccessorBlocks(), getPredecessorBlocks()`
- `getTrueBranch(), getFalseBranch()`
- `method`

Static Context

Provide relevant static information at joinpoints

■ Instruction

- `index, opcode, next, previous, isConditionalJump() ...`
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- `methodOwner, methodName, currentClassName`
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■ BasicBlock

- `id, index, blockType, size, getFirstInstruction() ...`
- `getSuccessorBlocks(), getPredecessorBlocks()`
- `getTrueBranch(), getFalseBranch()`
- `method`

■ Method

- `name, getNumberOfBlocks(), getEntryBlock(), getExitBlocks()`
- `node, classContext`

■ Class: `name, node`

Example – BasicBlock Execution

Intercepting basic block executions with BISM

```
public class BasicBlockTransformer extends Transformer {

    @Override
    public void onBasicBlockEnter(BasicBlock bb){
        String blockId =
            bb.method.className+"."+bb.method.name+"."+bb.id;

        print("Entered block:" + blockId)
    }

    @Override
    public void onBasicBlockExit(BasicBlock bb)
        String blockId =
            bb.method.className+"."+bb.method.name+"."+bb.id;

        print("Exited block:" + blockId)
    }
}
```

Dynamic Context

Provide access to dynamic values possibly only known during execution

■ Common

- `getThis()`
- `getLocalVariable(int index)`
- `getStackValue(int index)`
- `getInstanceField(String name)`
- `getStaticField(String name)`

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■ Adding new local variables

- `addLocalVariable(Object value)`
- `updateLocalVariable(LocalVariable, Object value)`

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BISM weaves necessary bytecode to extract values from stack or local variables

Instrumentation Methods

Methods to instrument target program to emit events or modify code

■ Printing

- `print(String)`, `print(DynamicValue)`, ...
- `printHash(DynamicValue)`, ...

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- `new StaticInvocation(..)` object should be constructed
- `addParameter()` to add primitive or `DynamicValue`
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■ Raw bytecode insertion

- `insert(AbstractInsnNode ins)`
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- `insert(AbstractInsnNode ins)`
- `insert(List<AbstractInsnNode> ins)`

Calls are compiled by BISM into bytecode instructions and inlined at needed locations

Example – Monitoring a List Iterator

```

public class IteratorTransformer extends Transformer {
    @Override
    public void afterMethodCall(MethodCall mc, MethodCallDynamicContext
        dc){
        if (mc.methodName.equals("iterator") &&
            mc.methodOwner.endsWith("List")) {
            // Access to dynamic data
            DynamicValue list = dc.getMethodTarget(mc);
            DynamicValue iterator = dc.getMethodResult(mc);

            // Invoking a monitor
            StaticInvocation sti =
                new StaticInvocation("IteratorMonitor",
                    "iteratorCreation");
            sti.addParameter(list);
            sti.addParameter(iterator);
            invoke(sti);
        }
    }
}

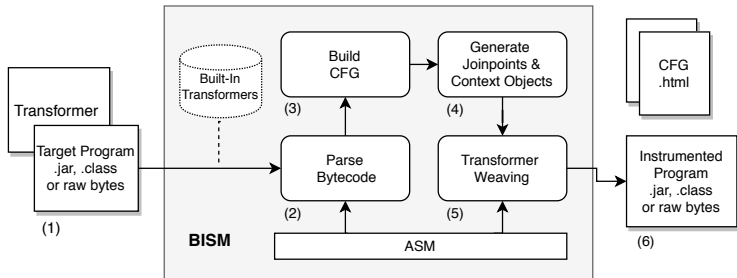
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BISM Implementation

- Open-source implemented in Java (\approx 4,000 LOC and 40 classes)
- Uses ASM for bytecode parsing, analysis, and weaving



Instrumentation process in BISM

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AES (Advanced Encryption Standard)

■ Experiment

- Instrument external AES implementation in build-time mode with BISM and DiSL
- Deploy inline monitors to report test inversions by duplicating all conditional jumps in successor basic blocks

```
// duplicating the
// if-statements
if(g_ptc > 0) {
    if(g_ptc > 0){
        ...
    }
    else {
        Report Attack!
    }
}
else{
    if(g_ptc > 0 ) {
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```

Needed Instrumentation

AES (Advanced Encryption Standard)

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- Instrument external AES implementation in build-time mode with BISM and DiSL
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■ Measurement

- Execution time and used memory on different input files
- Median of 100 runs for each input file

■ Experimental Setup

- Java JDK 8u181 with 4 GB maximum heap
- Intel Core i72.2 GHz, 16 GB RAM
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Needed Instrumentation

Inlining the monitor cannot be done with AspectJ

AES – DiSL Implementation

```

@SyntheticLocal
static int target = -1; //similarly op1,op2,opcode

@Before(marker = BytecodeMarker.class, args = "IF_ICMPGT,IF_ICMPLE..")
static void onIntegerCMP(InstructionStaticContextFull bcsc ..) {
    opcode = bcsc.getOpcode();
    op2 = dc.getStackValue(0, int.class);
    op1 = dc.getStackValue(1, int.class);
    target = bcsc.getJumpTargetLabel();
}

@Before(marker = BasicBlockMarker.class)
public static void afterbb(InstructionFullStaticContext bcsc ..) {
    boolean tt = target == ((int) bcsc.getIndexOfRegionStart());
    boolean attack = false;
    switch (opcode) {
        case 164: // IF_ICMPLE(164),
            if (((int) op1 <= (int) op2) != tt) {
                attack = true;
            }
            break;
    }
    ...

```

AES – BISM Implementation

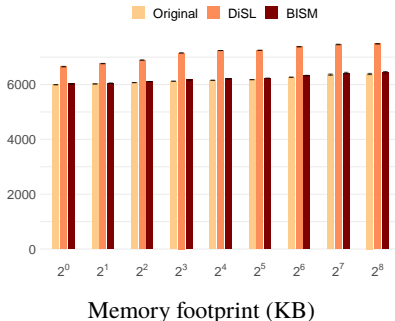
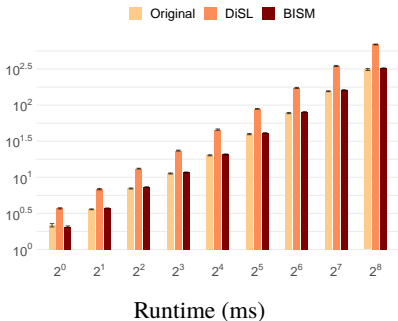
```
@Override
public void beforeInstruction(Instruction ins ..) {
    if (ins.isConditionalJump()) {
        if (ins.stackOperandsCountIfConditionalJump() == 1) {
            insert(new InsnNode(OpCodes.DUP));
        } // else use DUP2
    }

@Override
public void onTrueBranchEnter(BasicBlock bb ..) {
    LabelNode l_1 = new LabelNode();

    insert(new JumpInsnNode(bb.getLastRealInstruction().opcode, l_1));
    print("Test Inversion Attack!");
    insert(l_1);
}
```

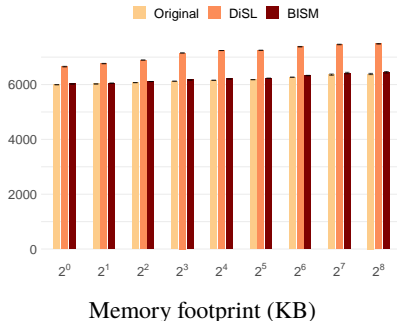
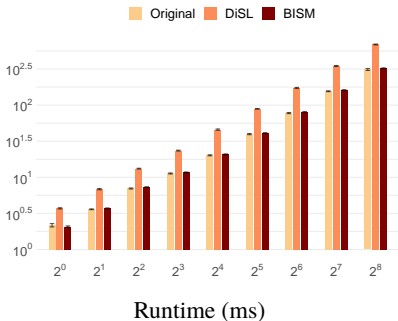
AES Results – BISM vs. DiSL

Input File (KB)	2 ⁰	2 ¹	2 ²	2 ³	2 ⁴	2 ⁵	2 ⁶	2 ⁷	2 ⁸
Events (M)	0.92	1.82	3.65	7.34	14.94	29.53	58.50	117.24	233.10



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Input File (KB)	2 ⁰	2 ¹	2 ²	2 ³	2 ⁴	2 ⁵	2 ⁶	2 ⁷	2 ⁸
Events (M)	0.92	1.82	3.65	7.34	14.94	29.53	58.50	117.24	233.10



**Less execution time and memory overhead than DiSL
for all input file sizes**

Financial Transaction System

■ Experiment

- Instrument an abstraction of a Financial Transaction System in build-time mode, with BISM and DiSL, to emit events `begin(i)` and `end(i)`
- Deploy outline monitor to report arbitrary jumps

```
// emit events begin(i)
    & end(i)

// Basic Block i
begin(i)
first instruction
    ...
last instruction
end(i)

// begin(i) & end(i) are
    function calls
```

Needed Instrumentation

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- Execution time and used memory on two scenarios
- Median of 100 runs for each scenario

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Needed Instrumentation

Transactions – DiSL Implementation

```
public class BasicBlockExtendedContext extends BasicBlockStaticContext {
    public String uniqueBlockId() {
        MethodNode method = staticContextData.getMethodNode();
        String uniqueId = staticContextData.getClassNode().name
            + "." + method.name + "." + method.signature + "." + this.getIndex ();
        return uniqueId;
    }
}
```

```
@Before(marker = BasicBlockMarker.class, scope = "transaction.*.*")
public static void beforebbb( BasicBlockExtendedContext bc,
    MethodStaticContext mc ) {
    transaction.MonitorAJ.begin(bc.uniqueBlockId());
}

@After(marker = BasicBlockMarker.class, scope = "transaction.*.*" )
public static void afterbbb( BasicBlockExtendedContext bc,
    MethodStaticContext mc) {
    transaction.MonitorAJ.end(bc.uniqueBlockId());
}
```

Transactions – BISM Implementation

```
@Override
public void onBasicBlockEnter(BasicBlock bb){
    String uniqueBlockId = bb.method.className+"."+bb.method.name+"."
        +bb.method.methodNode.signature+"."+bb.index;

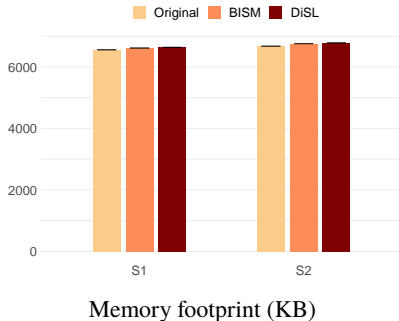
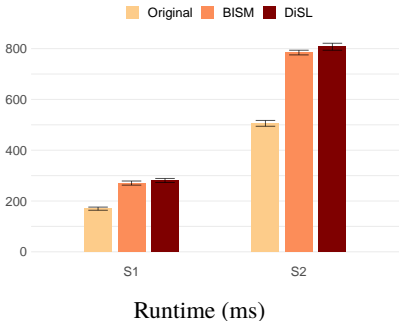
    StaticInvocation sti =
        new StaticInvocation(monitor, "begin");
    sti.addParameter(uniqueBlockId);
    invoke(sti);
}

@Override
public void onBasicBlockExit(BasicBlock bb){
    String uniqueBlockId = bb.method.className+"."+bb.method.name+"."
        +bb.method.methodNode.signature+"."+bb.index;

    StaticInvocation sti =
        new StaticInvocation(monitor, "end");
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}
```

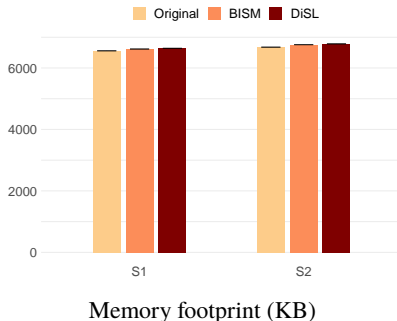
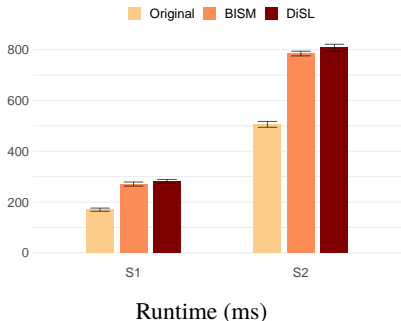
Transactions Results – BISM vs. DiSL

Scenario	S1	S2
Events (M)	679	2307



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**Less execution time and memory overhead than DiSL
for both scenarios**

DaCapo Benchmarks

We compare BISM, DiSL & AspectJ in a general RV scenario

■ Experiment

- Use **HasNext**, **UnsafeIterator** and **SafeSyncMap** properties
- External monitor library with stub methods to receive extracted objects
- **DaCapo** suite (dacapo-9.12-bach) targeting only packages specific each benchmark

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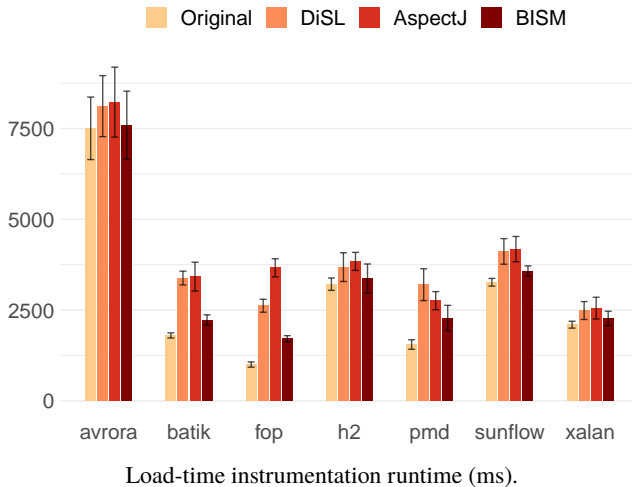
■ Measurement

- **Two modes: Build-time** and **Load-time** instrumentation
- Median of 100 runs for each benchmark for each mode

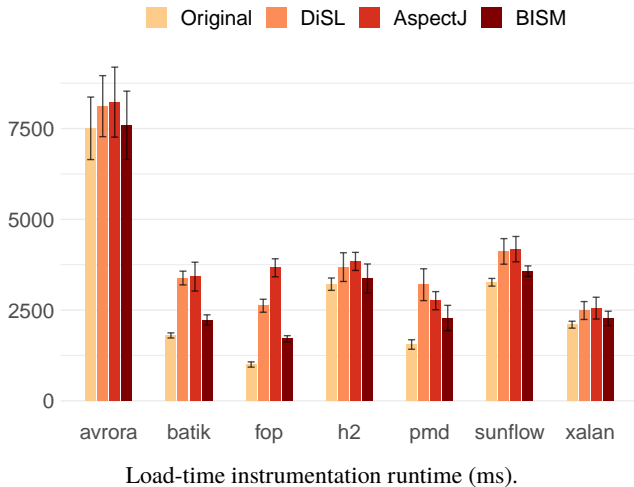
■ Experimental Setup

- Java JDK 8u251 with 2 GB maximum heap size
- Intel Core i9-9980HK (2.4 GHz. 8 GB RAM)
- Ubuntu 20.04 LTS 64-bit

DaCapo – Load-time Results

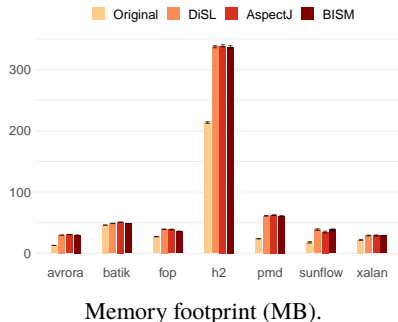
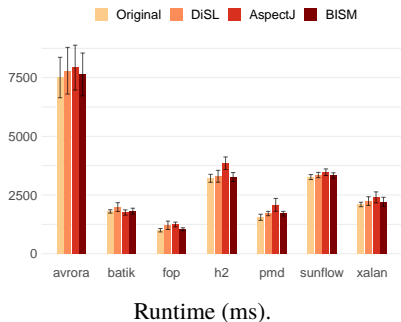


DaCapo – Load-time Results

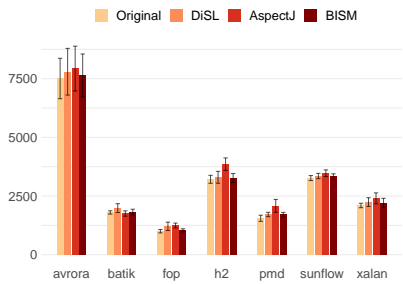


BISM shows better performance over DiSL and AspectJ in all benchmarks

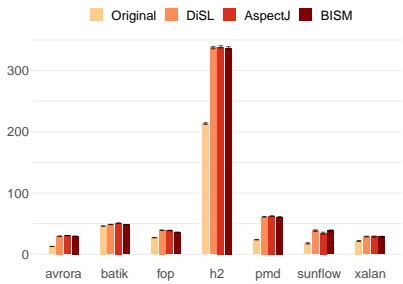
DaCapo – Build-time Results



DaCapo – Build-time Results



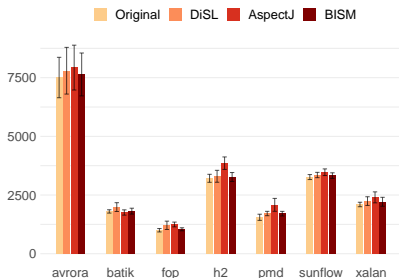
Runtime (ms).



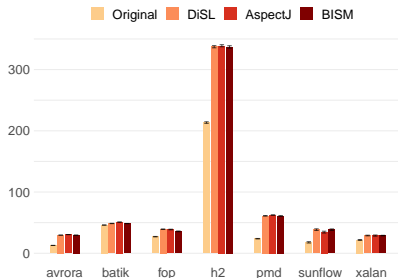
Memory footprint (MB).

- BISM shows less overhead in all benchmarks in execution time, except for batik, where AspectJ emits fewer events

DaCapo – Build-time Results



Runtime (ms).



Memory footprint (MB).

- BISM shows less overhead in all benchmarks in execution time, except for batik, where AspectJ emits fewer events
- BISM also shows less overhead in used-memory footprint, except for sunflow, where AspectJ emits much fewer events

DaCapo – More Results

	Scope	Instr.	BISM	DiSL	AspectJ	Events (M)	
			Ovh.%	Ovh.%	Ovh.%	#	AJ
avroa	1,550	35	2.72	5.06	34.24	2.5	2.5
batik	2,689	136	1.81	2.85	9.59	0.5	0.4
fop	1,336	172	1.35	5.16	27.07	1.6	1.5
h2	472	61	1.44	3.75	37.75	28	28
pmd	721	90	2.38	5.03	29.63	6.6	6.3
sunflow	221	8	2.90	7.25	23.19	3.9	2.6
xalan	661	9	1.00	3.00	16.00	1	1

Bytecode size of the instrumented benchmarks applications

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Bytecode size of the instrumented benchmarks applications

BISM incurs less bytecode size overhead for all benchmarks

Outline

- 1 Motivation
- 2 BISM Design Goals & Features
- 3 BISM Language
- 4 BISM Implementation
- 5 Evaluation
- 6 Conclusion**

Conclusion

- BISM is effective for low-level and control-flow aware instrumentation and can be used for lightweight and expressive runtime verification
- BISM shows better performance than BISM and AspectJ

	Overhead reduction	
	DiSL	AspectJ
Bytecode Performance		
Execution time	45%	82%
Used Memory	5%	8%
Size	56%	91%
Tool performance		
Execution time	64%	68%

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	BCEL	ASM	Javassist	CGLIB	DiSL	AspectJ	BISM
Bytecode Visibility	✓	✓	✓	✓	✓	✗	✓
Allows bytecode insertion	✓	✓	✓	✓	✗	✗	✓
No bytecode proficiency	✗	✗	✗	✗	✓	✓	✓
High level of abstraction	✗	✗	✗	✗	✓	✓	✓
Follows AOP paradigm	✗	✗	✗	✗	✓	✓	✓

The Future of BISM

We plan on adding

- More instrumentation methods like calling non-static external methods
- More built-in Transformers covering classic RV use cases
- Instrumentation directives
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New directions now possible with BISM

- Control-flow-aware RV
- Runtime Enforcement using bytecode insertion capabilities