PROJET COFFI
ANR-18-CES39-003

VERIFICATION OF THE PROGRAM'S CFI BASED ON A TRACE ENCODER

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Journée thématique sur les attaques par injection de fautes

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INSPIRING INNOVATION SINCE 1816

Table of Contents

Table of Contents

Introduction

Why is it necessary to guarantee the program's CFI?

- Software attacks : Buffer Overflow, ROP, Code Reutilization Attacks...
- Hardware attacks : Fault Injection...
- Example - Code pin verification :

```
   int counter = 3;
   void VerifyPin() {
1 →   if (counter > 0)
2 →     if (Cmp(userPIN,devicePIN))
          Accept();
        else
3 →       counter--;
   }
```

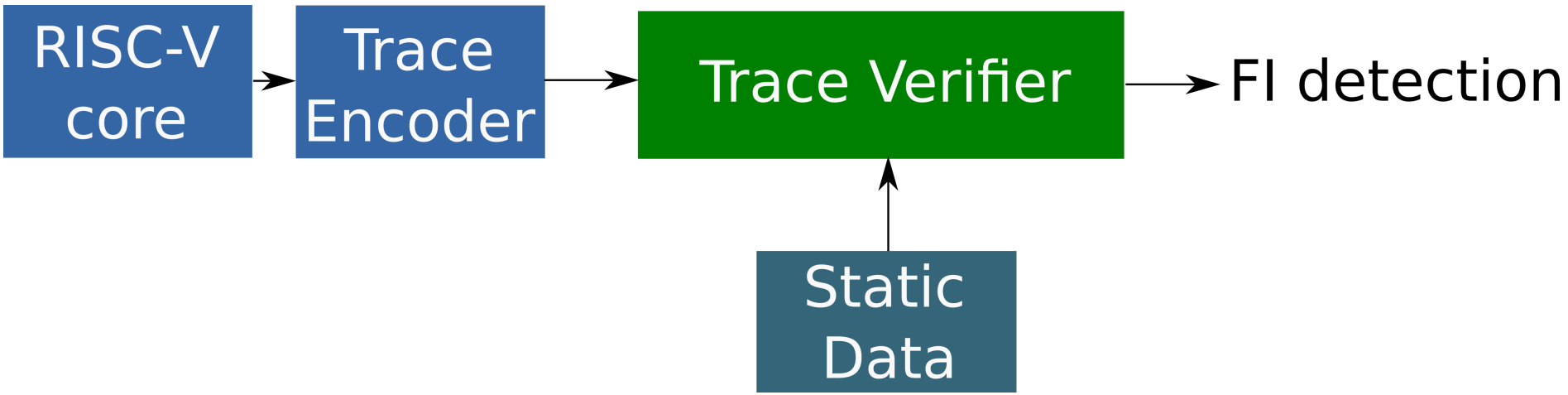
- CFI Verification = Correct Program Execution.

Table of Contents

How could we verify the program's CFI?

Methodology

- Get information about what is executed at the RISC-V core level => Trace Encoder (TE)
- Compare these data to static data obtained from a static analysis of the binary program => Trace Verifier (TV)
- Detect if a fault injection attack is made => TV's output

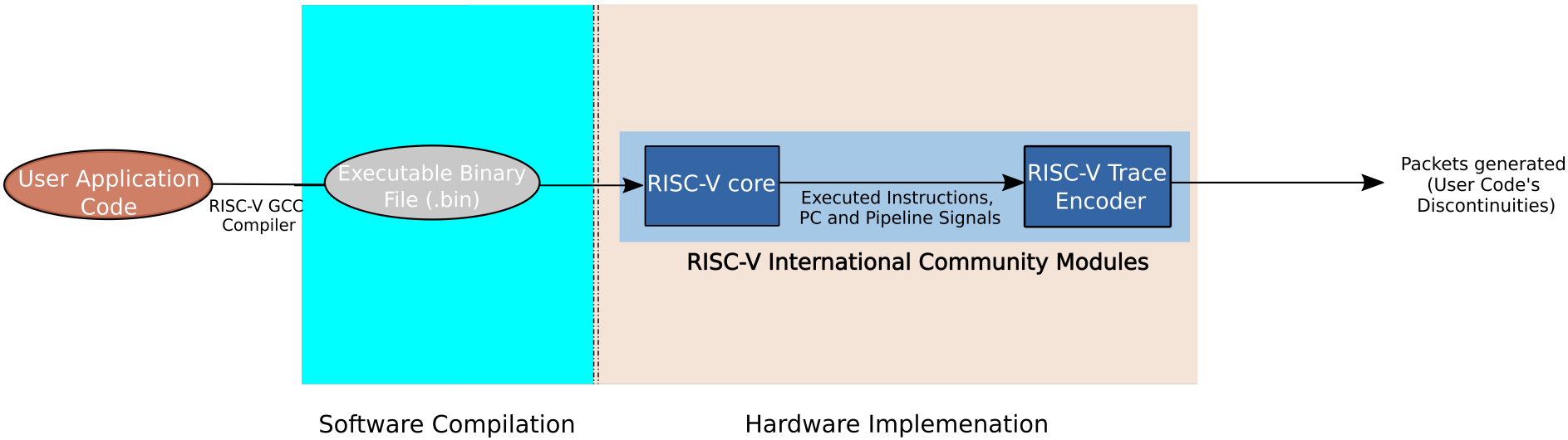


Trace Encoder

Definition

- Module designed by the RISC-V community.
- Overall objective: Compression of the program's execution path.
- Interpret the executed instructions from the RISC-V core.
- Report the discontinuities present in a program in the form of packets.
  - Instructions presenting an uninferable PC discontinuity.
  - Interruptions and exceptions...

Packet generation example



Example of a packet - Use Case

- Having a function call where the program encounter n branches.
- A packet will be sent containing : the number of branches (n), the branch map (branch taken or not) and the return address ...

Packet generation example

```
int fnct1(int a, int b, int c) {
        if (a=1) {
                ...
                if(b=2) {
                        ...
                } else {
                        ...
                }
                end if;
        } else {
                ...
        }
        if(c=a+b) {
                ...
        } else {
                ...
        }
return a+b+c;
}

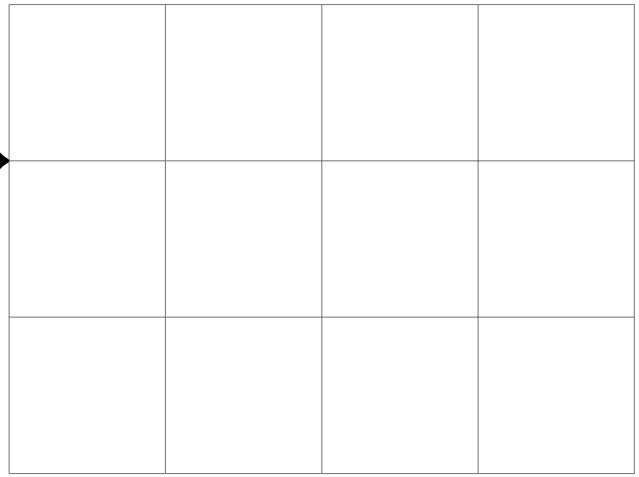
int main() {

int a = 1;
int b = 3;
int c = 4;

c=fnct1(a,b,c);

return 0;
}
```

Fnct1 call →



Packet generation example

```
1 →  int fnct1(int a, int b, int c) {
            if (a=1) {
                    ...
                    if(b=2) {
                            ...
                    } else {
                            ...
                    }
                    end if;
            } else {
                    ...
            }
            if(c=a+b) {
                    ...
            } else {
                    ...
            }
    return a+b+c;
    }

    int main() {

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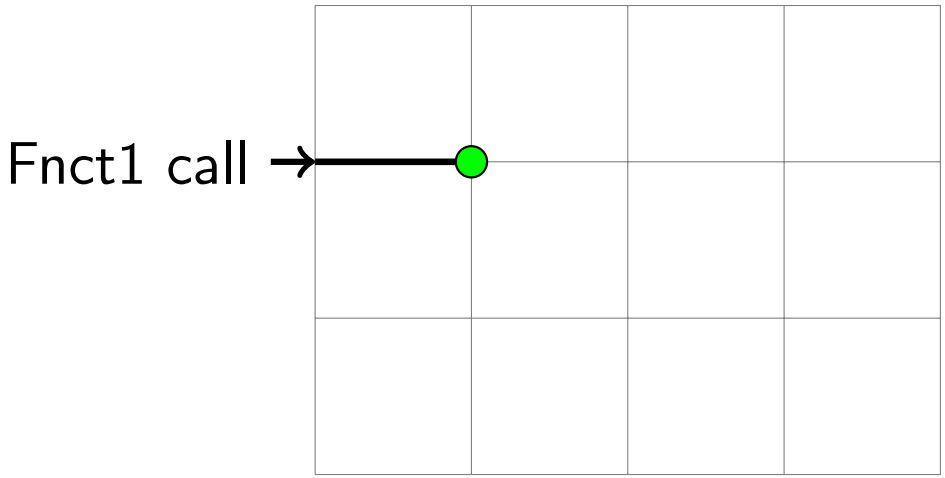

Packet generation example

```
1 →     int fnct1(int a, int b, int c) {
        if (a=1) {
2 →             ...
                if(b=2) {
                        ...
                } else {
                        ...
                }
                end if;
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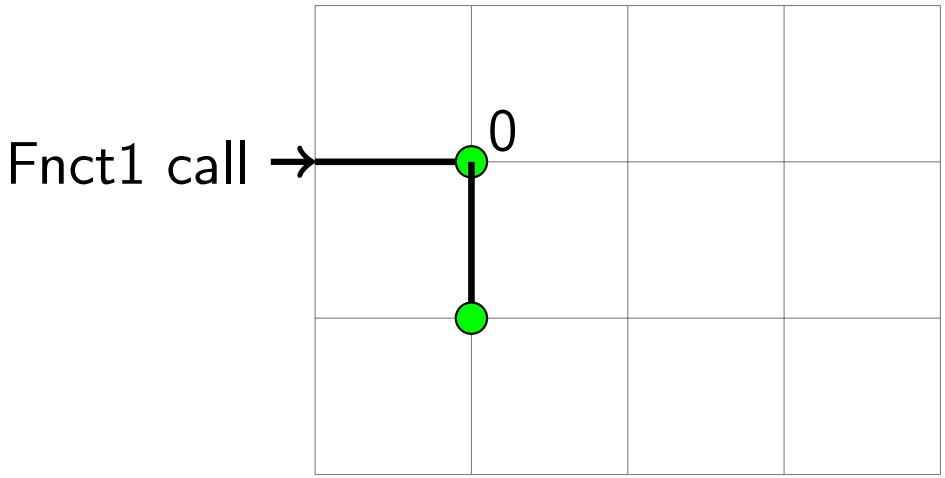

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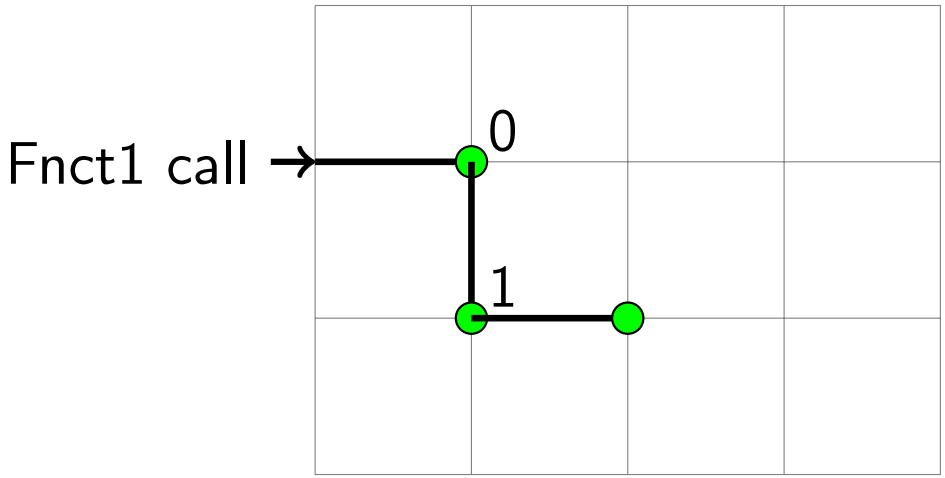

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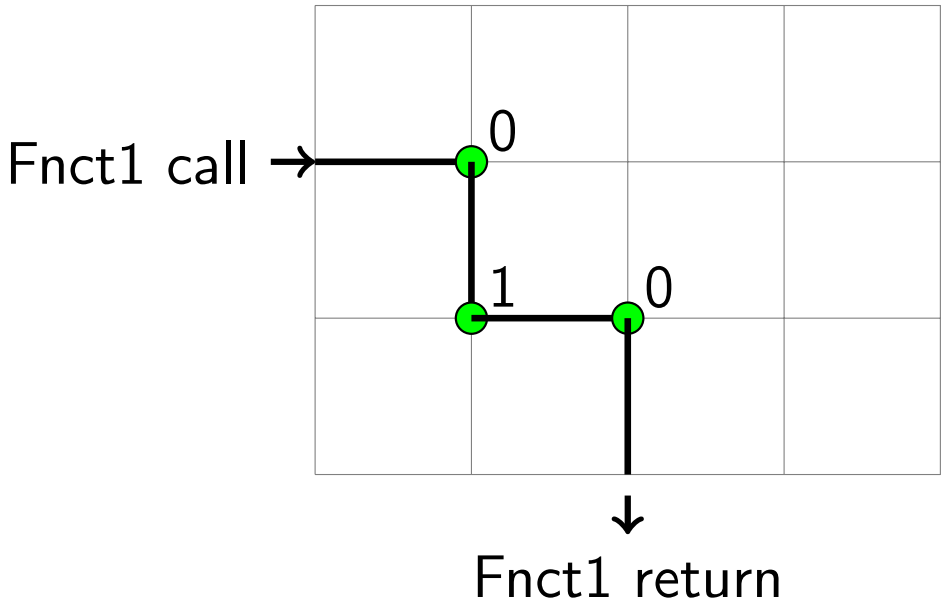

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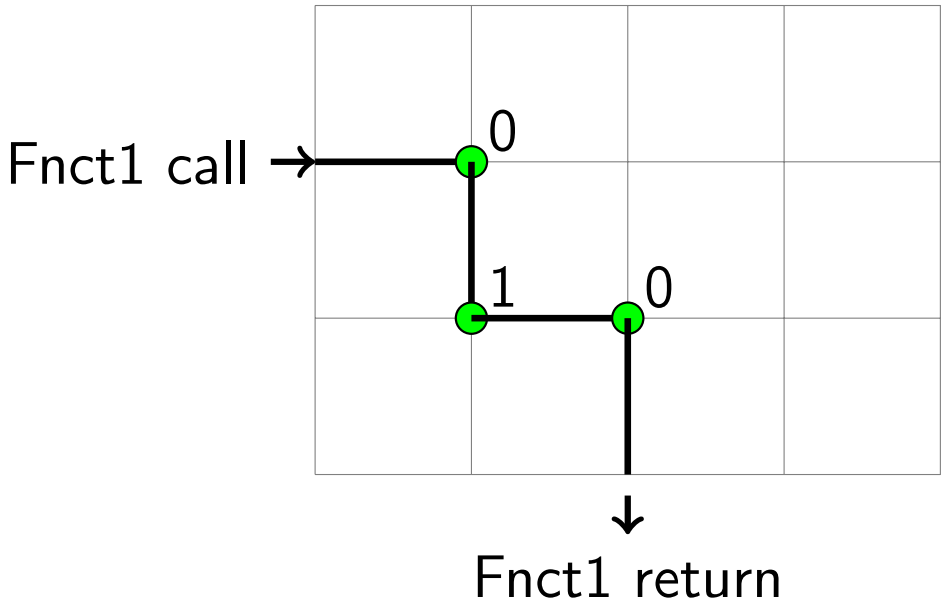
```

1 →
2 →
3 →



Example (Generated packet)

- Number of branches : 3
- Branch Map : 010
- Return address

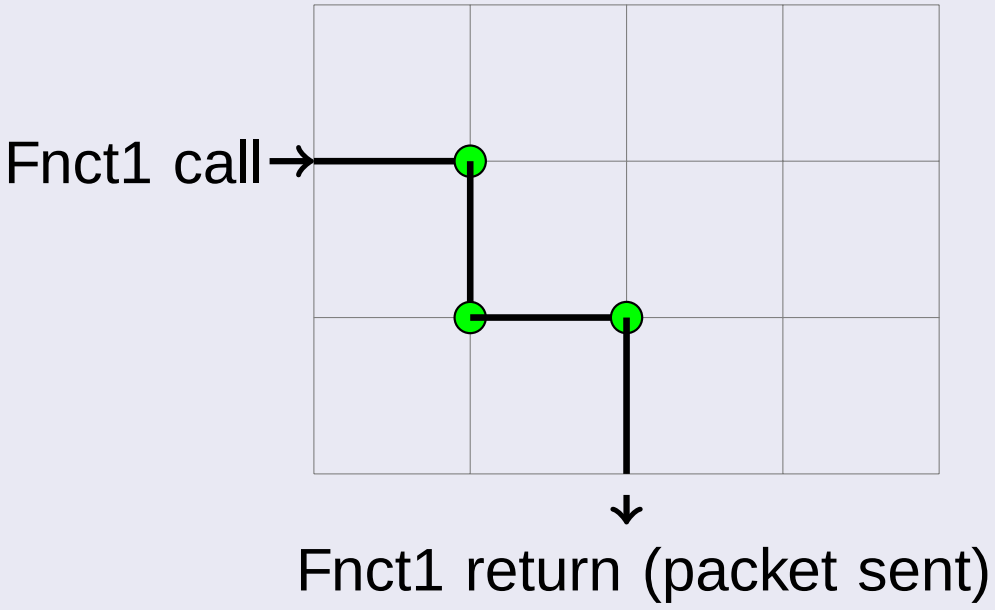
Trace Verifier

- It's verification system based on the TE sent packets.
- Compare the TE packets to static data issued from a static analysis of the binary program.
  - A static analysis is made after the compilation process.
  - Branch, jump, call and return instructions with their addresses are stored in a memory.
- A flag is raised when a fault is detected.

Table of Contents

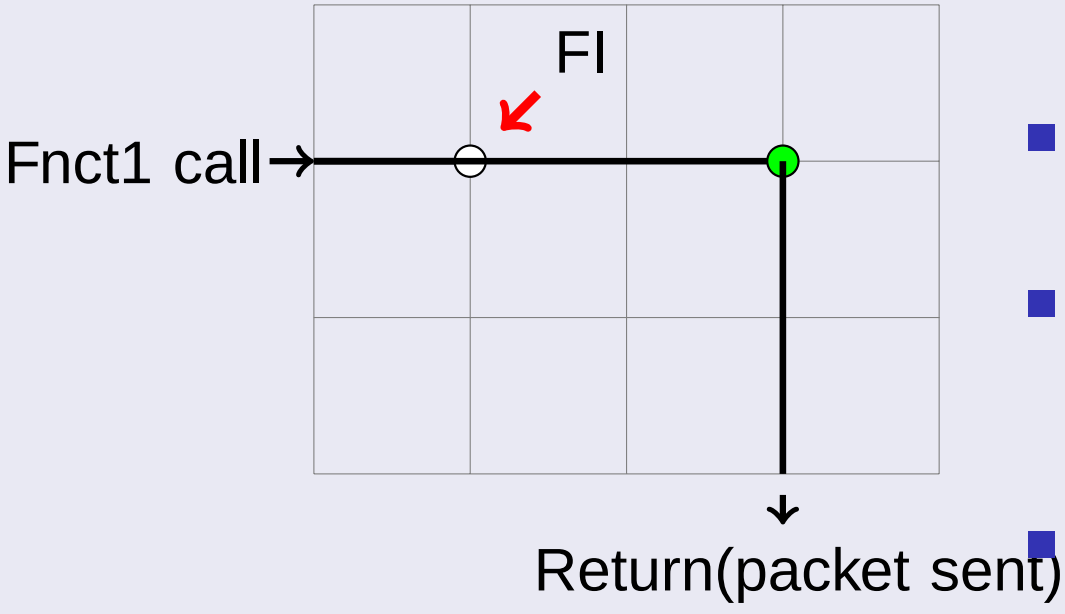
Exploited solution - Architecture

First approach - ASIC RISC-V Model



- Veri cation process starts when a packet is sent.
- Navigation through the static data and constitution of the path followed by the program.
- Comparison process.

First approach - ASIC RISC-V Model



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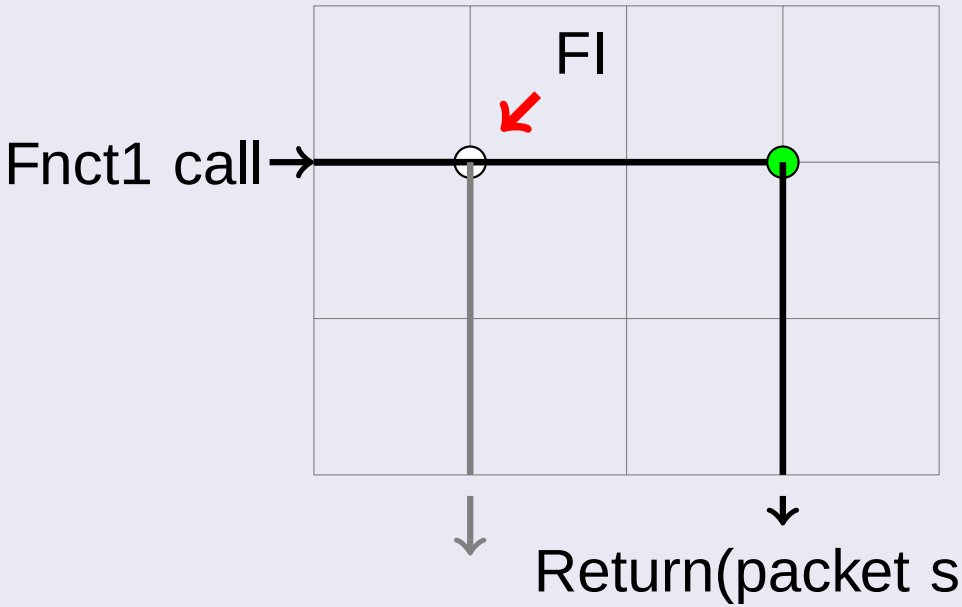

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- Comparison process.

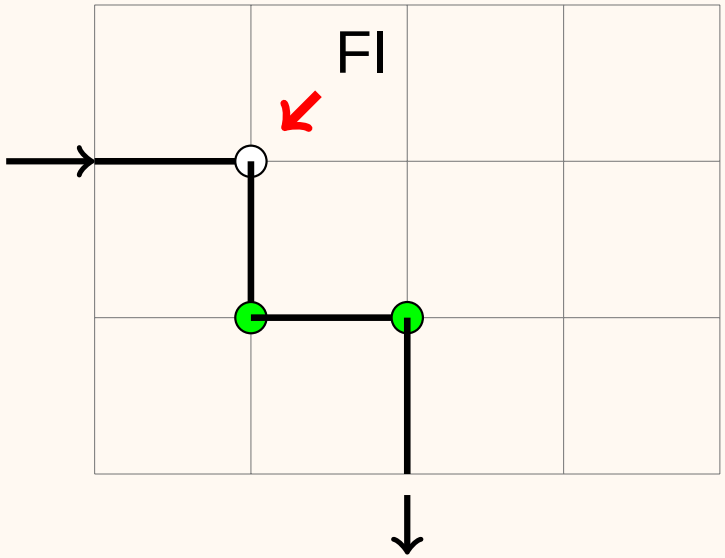
First approach - FI detection

Fault coverage

- Changing the return address of a function.
- Instruction skip on a function call.

Particular cases

- Instruction skip on branch instructions.
- It depends on the branch number and return address.



Limitations

- Corruption of a branch instruction (funct3, branch address ...).

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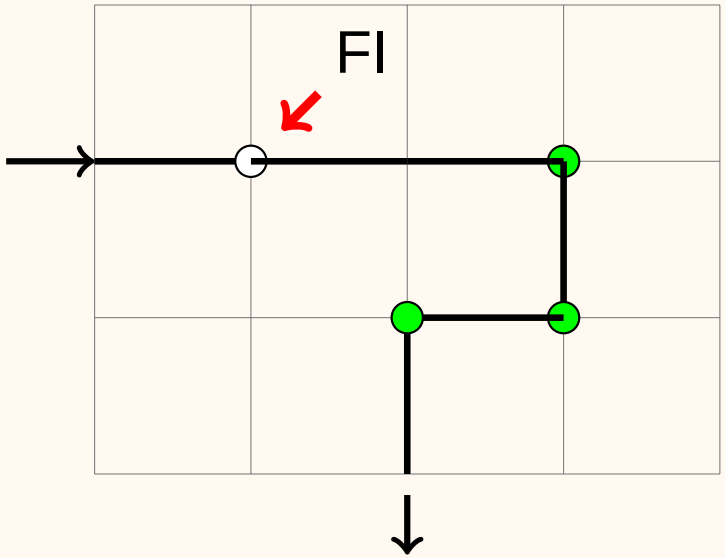

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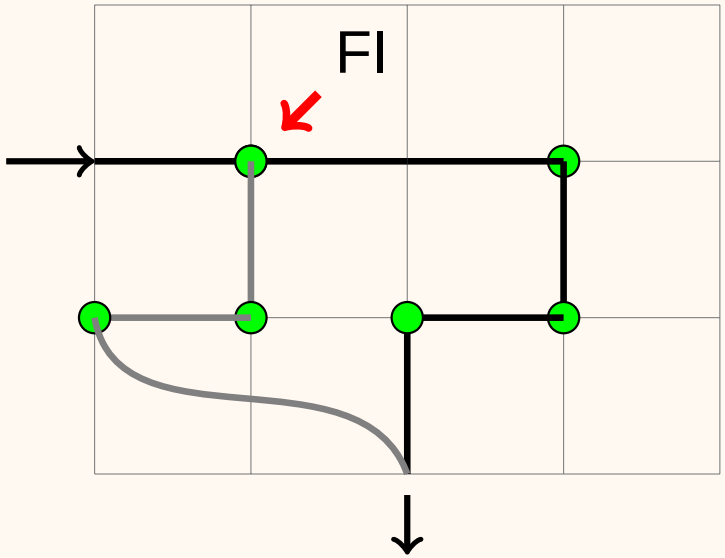


Limitations

- Corruption of a branch instruction (funct3, branch address ...).

Second approach - Adding the PC - RISC-V FPGA Mod

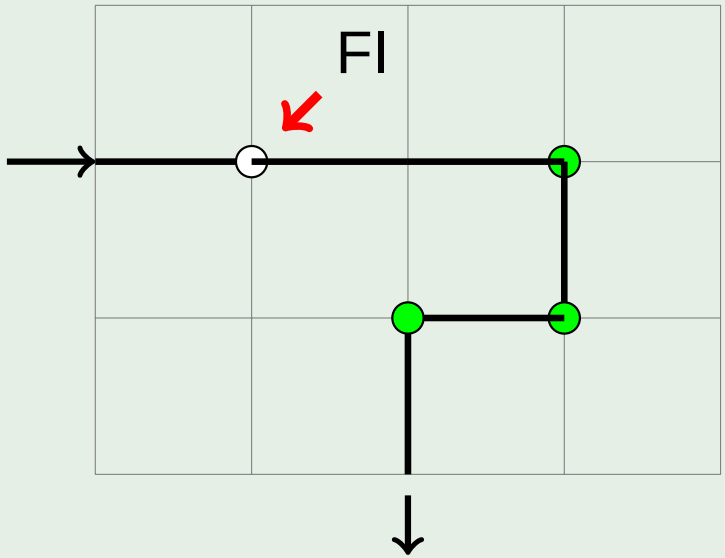
Features

- We pull the PC and connect it to the TV.
- Calculation is made before receiving the packet.
- Veri cation process is faster compared to the rst approach.

Second approach - FI detection

Fault coverage

- Changing the return address of a function.
- Instruction skip on a function call andbranch instruction .
  - Total number of branches=4.



Particular cases

- Changing the branch address in case it was taken : beq a4,a5, 1300

Limitations

- Corruption of a branch instruction (funct3, branch address ...).

Third approach - Rede ning the TE RISC-V standard

Features

- PC and executed instructions are pulled and connected to the TV.
- Adjustment to the TE RISC-V standard by de ning the quali ed instructions (jump, branch, return...).
- Packet is sent after each quali ed instruction.

Third approach - FI detection

Fault coverage

- Changing the return address of a call function.
- Instruction skip on a function call and branch instruction.
- Their corruption / substitution with other instructions.

Table of Contents

Results - Comparison between the 3 approaches

- Each approach covers a speci c number of threats :

Approach	SFC	RAC	SBI	CDI	L
TV-ASIC	✓	✓	(X)	X	- -
TV-FPGA-PC	✓	✓	✓	X	-
TE-TV-CFI	✓	✓	✓	✓	+

- SFC : Skip on function calls.
- RAC : Return address change.
- SBI : Skip on branch instructions.
- CDI : Corruption of a discontinuity instruction.
- L : Latency.

Conclusion - Comparison between the 3 approaches

- Hardware Area Overhead (in terms of slices) :

Approach	TE	TV	Total
TV-ASIC	241	360	601
TV-FPGA-PC	241	641	882
TE-TV-CFI	95	575	670

- RISC-V IBEX : 635 slices.

Table of Contents

Perspectives

- De nition of a new packet for the basic blocks' veri cation, by adjusting the TE's standard.
- Veri cation of the correct instructions execution in the processor pipeline (cf. COFFI Project).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