

SAMVA: Static Analysis for Multi-Fault Attack Paths Determination

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Fault injection attacks

Adversary goals

- Leak critical data
- Break cryptographic properties
- Take over a device

Fault injection for control-flow hijacking

- Execute authentication code
- Avoid countermeasures

```
BOOL verifyPIN() {  
    g_authenticated = 0;  
  
    if(g_ptc > 0) {  
        if(byteArrayCompare(...) == 1) {  
            g_ptc = 3;  
            g_authenticated = 1;  
            return 1;  
        } else {  
            g_ptc--;  
            return 0;  
        }  
    }  
  
    return 0;  
}
```

Example with a source code in C



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Fault injection attacks

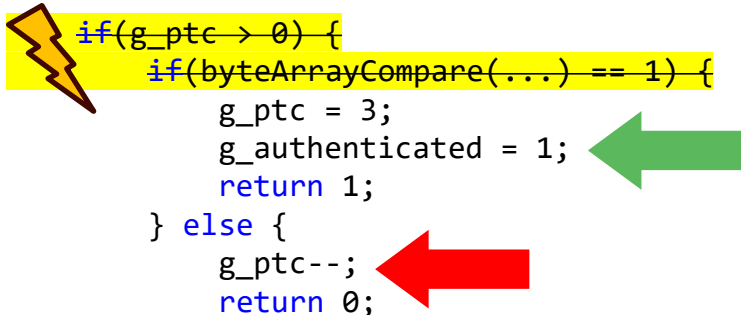
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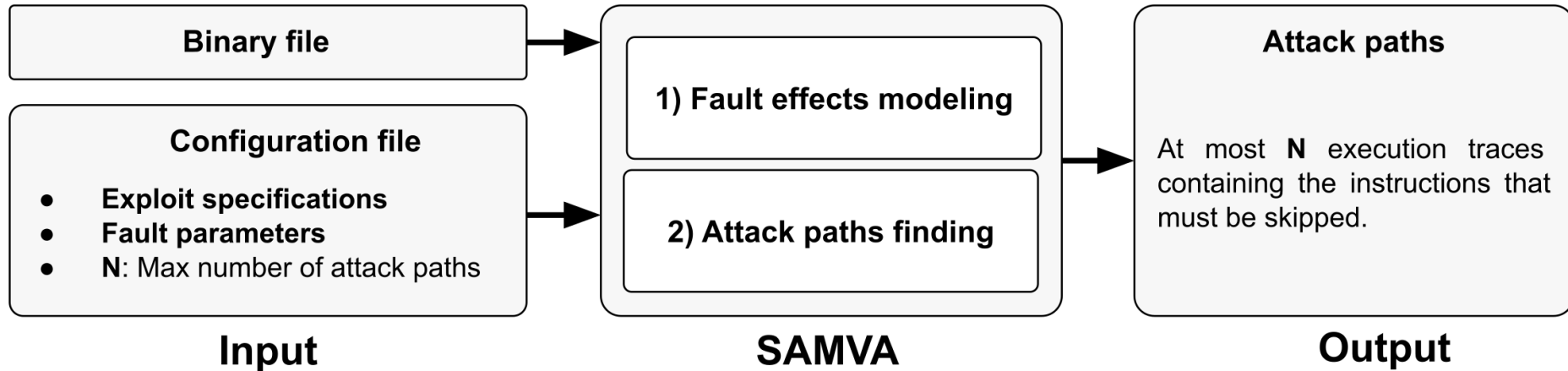
**How to assess the robustness
of a software against fault injection?**

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Example with a source code in C

Contribution: SAMVA

- **Static analysis method:** Finding attack paths semi-automatically
- **Multiple instruction-skip fault model:** Faults with a variable width
- **Accessibility exploit model:** Reach and avoid specified regions of code binary



Fault model: multiple instruction-skip

Fault Parameters

B2:

0x10510: str r0, [fp, #-12]

0x10514: ldr r3, [fp, #-12]

0x10518: cmp r3, #4

0x1051c: bgt 10530

B3:

0x10520: ldr r3, [fp, #-8]

0x10524: add r3, r3, #1

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0x1052c: b 1053c

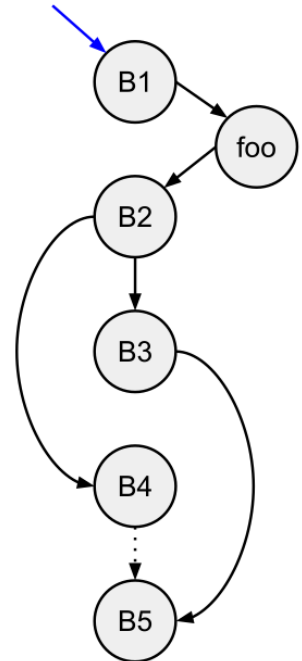
B4:

0x10530: ldr r3, [fp, #-8]

0x10534: sub r3, r3, #1

0x10538: str r3, [fp, #-8]

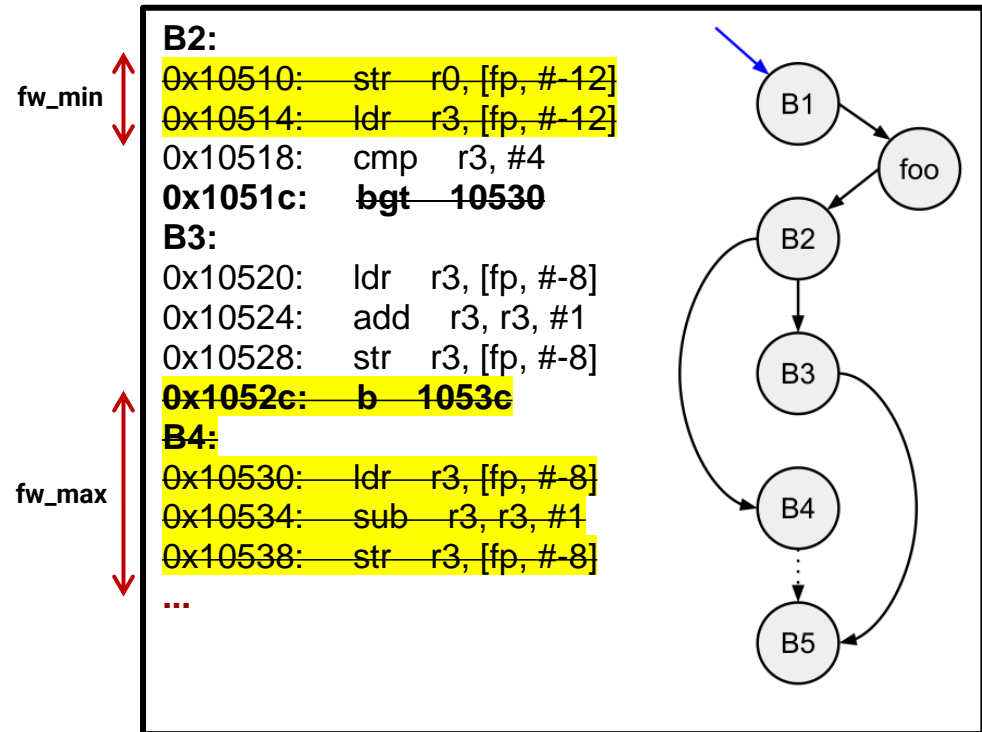
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Fault model: multiple instruction-skip

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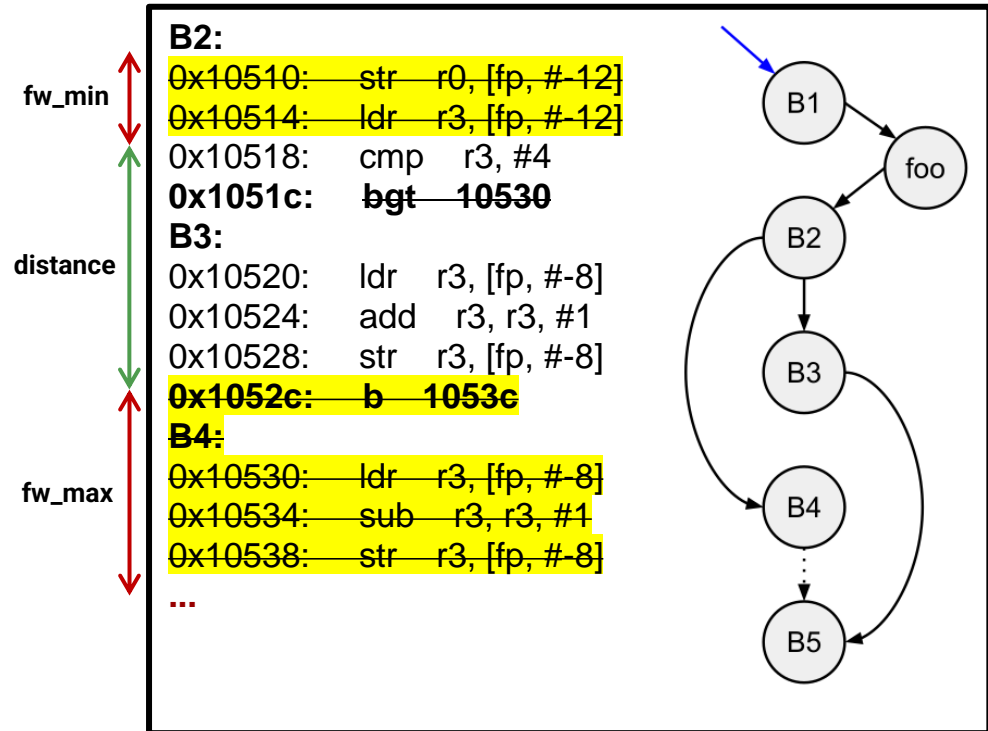
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Fault effects

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Predictable path w.o. data-flow analysis

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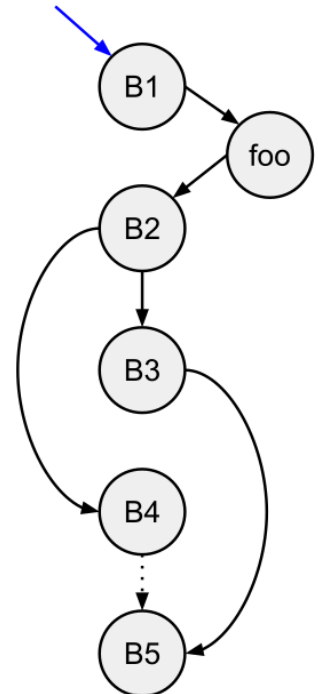
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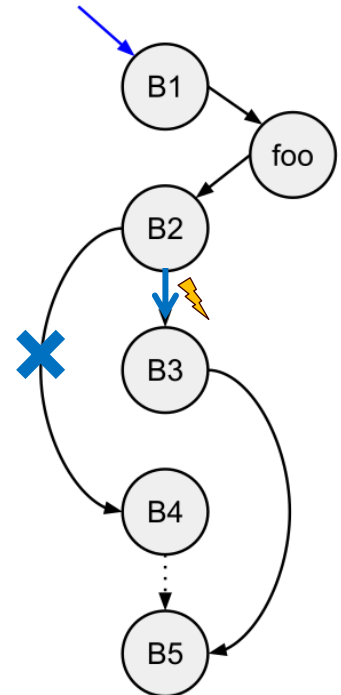
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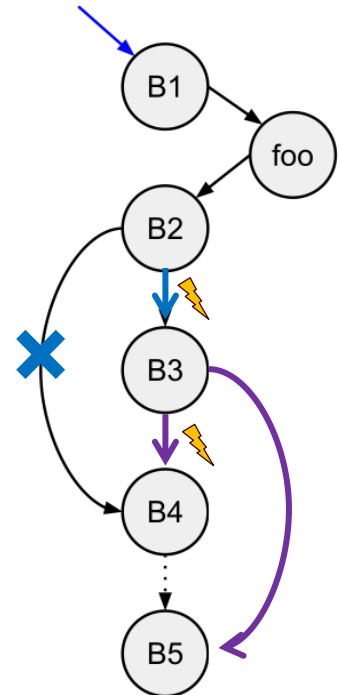
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Reflect these effects on the CFG

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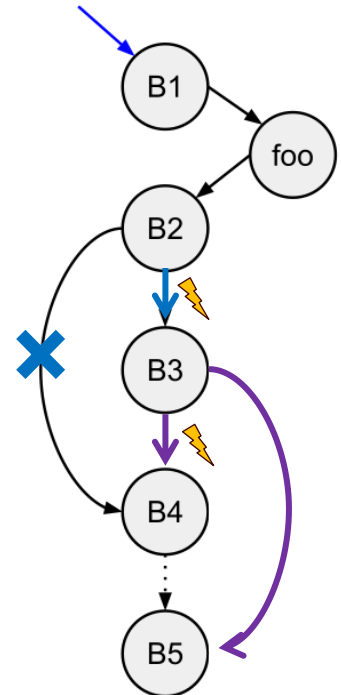
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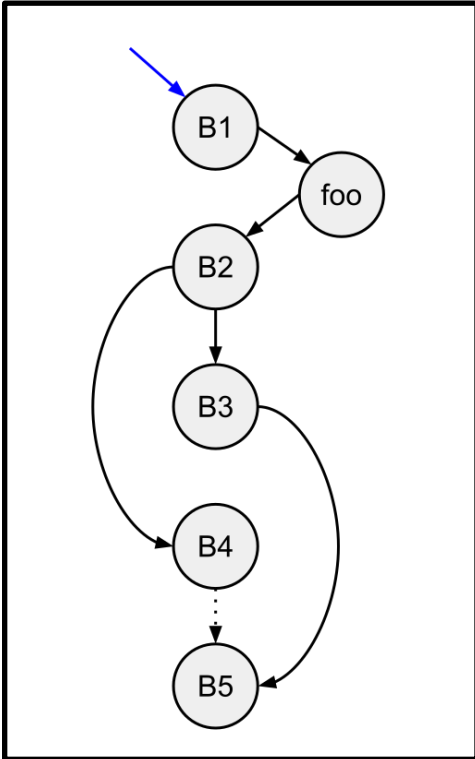
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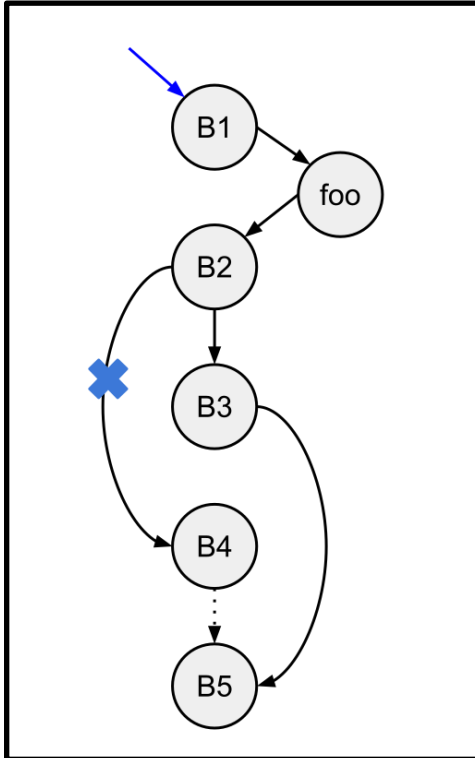


Fault effects modeling



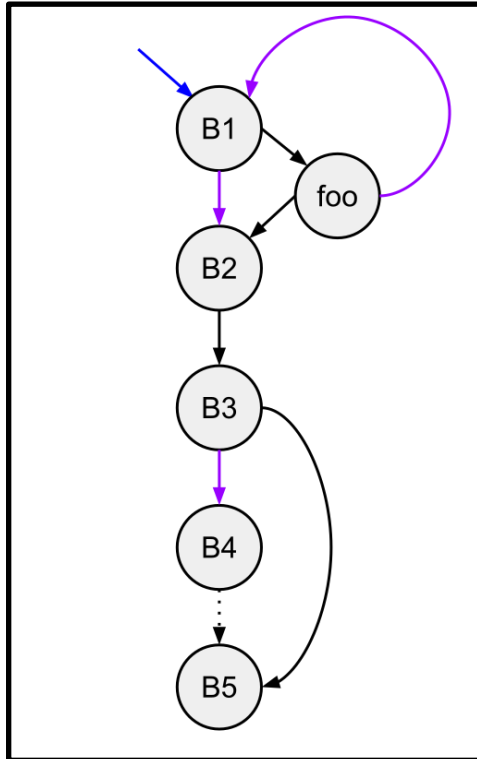
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Fault effects modeling



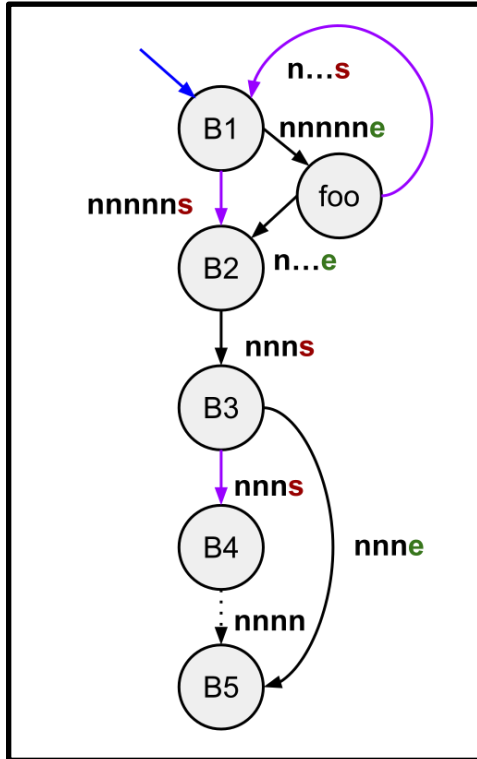
- 1) Retrieve the initial control flow graph (CFG)
- 2) Remove edges: BB ending with conditional jump

Fault effects modeling



- 1) Retrieve the initial control flow graph (CFG)
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- 3) Add new edges: BB ending with unconditional jump

Fault effects modeling



1) Retrieve the initial control flow graph (CFG)

2) Remove edges: BB ending with conditional jump

3) Add new edges: BB ending with unconditional jump

4) Annotate edges with a sequence of “types”

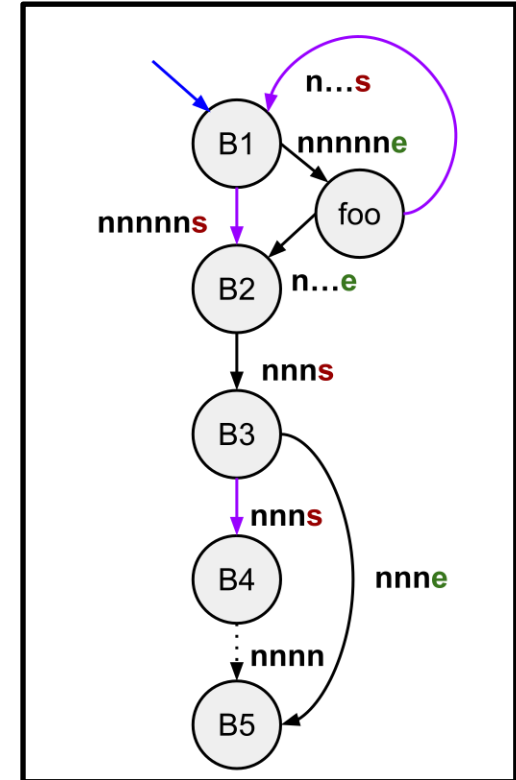
- One type per instruction of the source BB
 - execute (e)**: must be executed
 - skip (s)**: must be skipped
 - neutral (n)**: can be either skipped or executed

Attack paths finding

- Build a set of candidate paths

Example: Targeted basic blocks = [foo; B4]

B1 – foo – B2 – B3 – B4



Attack paths finding

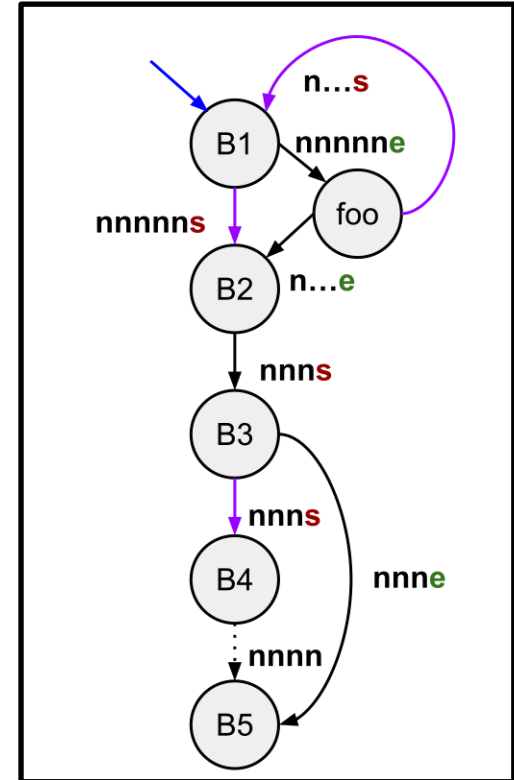
- **Build a set of candidate paths**

Example: Targeted basic blocks = [foo; B4]

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- **Build an execution trace**

- List of tuples <address, type>
- nnnnnne + nnnnnnnne + nnnns + nnnS



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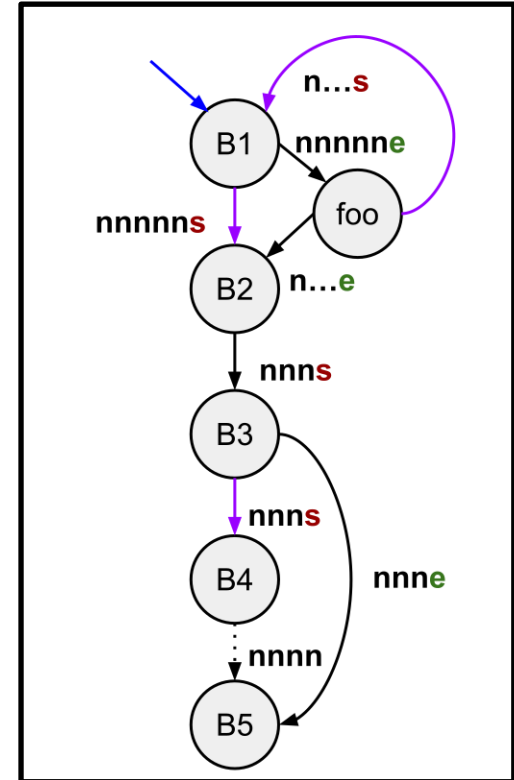
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- **Build an execution trace**

- List of tuples <address, type>
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- **Fault positioning**

Determine N set of faults {(position, width)}
making the path feasible



Fault positioning algorithm

0x104F8:	execute
0x104FC:	neutral
0x10500:	neutral
0x10504:	skip
0x10508:	skip
0x1050C:	neutral
0x10510:	execute
0x10514:	neutral
0x10518:	neutral
0x1051C:	skip
0x10520:	skip
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Execution trace example
 ($fw_min = 4$, $fw_max = 8$
 $fw_min_dist = 4$)

- **Conditions** for set of faults (a.k.a solution) to be valid
 - All instructions typed **skip** are covered by a fault
 - No instruction typed **execute** is covered by a fault
 - Faults widths $\in [fw_min, fw_max]$
 - Distances between two faults $\geq fw_min_dist$

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- **Conditions** for set of faults (a.k.a solution) to be valid
 - All instructions typed **skip** are covered by a fault
 - No instruction typed **execute** is covered by a fault
 - Faults widths $\in [fw_min, fw_max]$
 - Distances between two faults $\geq fw_min_dist$
- Solutions are built **incrementally** with a **backtracking** approach
- At the end, we obtain **N execution traces along with the position and the width of the faults for each**

Experimentation: Objective & Setup

Experimentation objective

Ensure that the attack paths found by SAMVA are really effective



Modified version of gem5 allowing the simulation of instruction-skips

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Modified version of gem5 allowing the simulation of instruction-skips

- **Benchmark: VerifyPIN suite**
 - FISCC [Dureuil et al. SAFECOMP 2016]
 - 8 implementations with increasing level of countermeasures
- **Tested with a large set of fault parameters**
 - **3366 different faults parameters** per binary
 - At most **30 attack paths per fault parameters**
 - Simulate attacks until one succeed

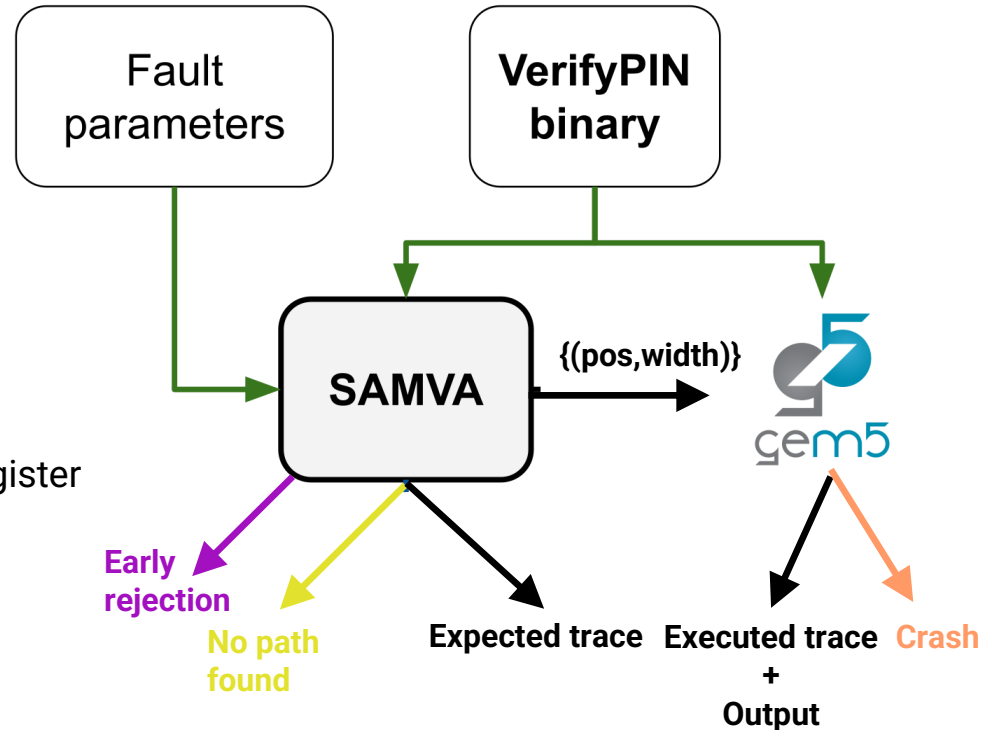
Validation methodology

Before Simulation

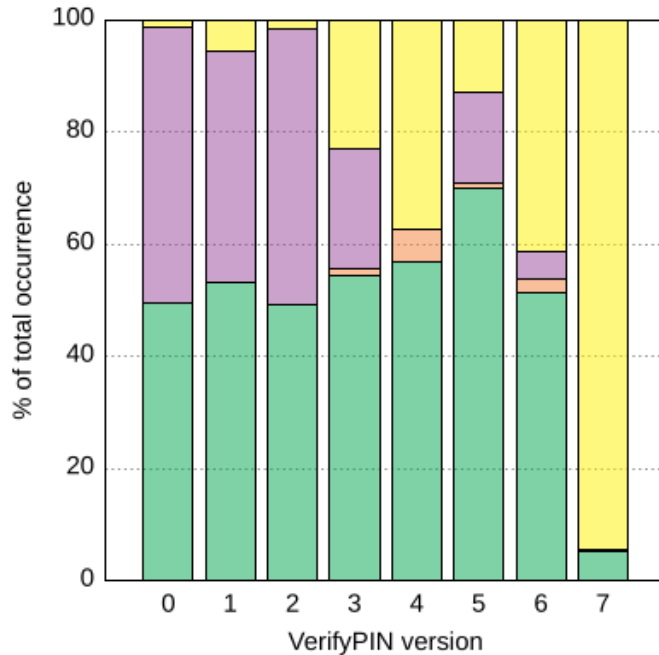
- **Early rejection**
Fault parameters not suited for the binary
- **No Path found**

After Simulation

- **Execution crashes due to faults**
e.g. Illegal load from address stored in a register
- **Validated attack**
Authenticated & Traces are equal



Classification of attack path searches

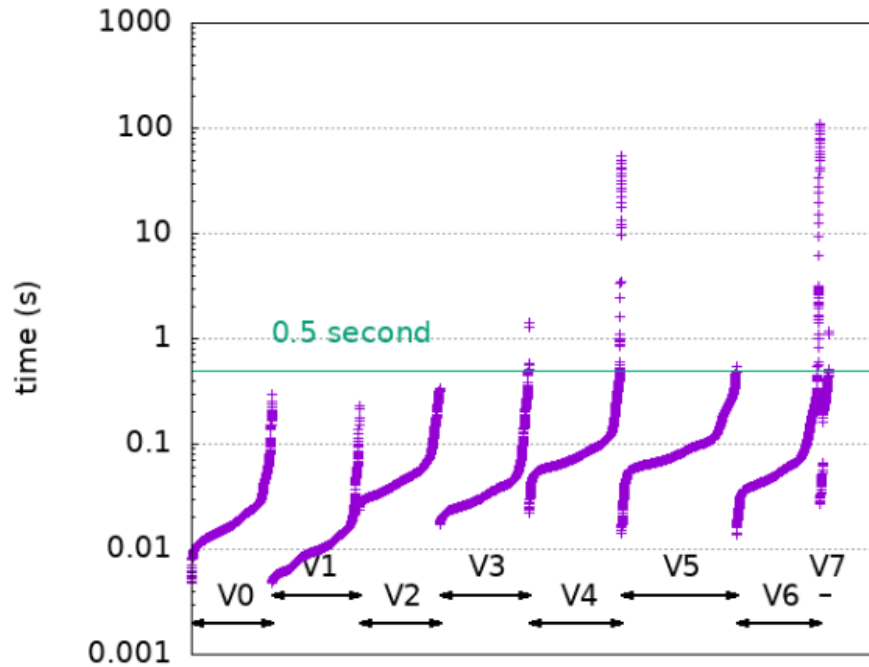


- All versions are vulnerable
- Versions **V4** and **V7** are the most robust implementations
- Remark: Facilitate fault positioning by $\searrow fw_min$, $\nearrow fw_max$ and $\searrow fw_min_dist$

→ **SAMVA is able to find numerous attacks paths for all versions**



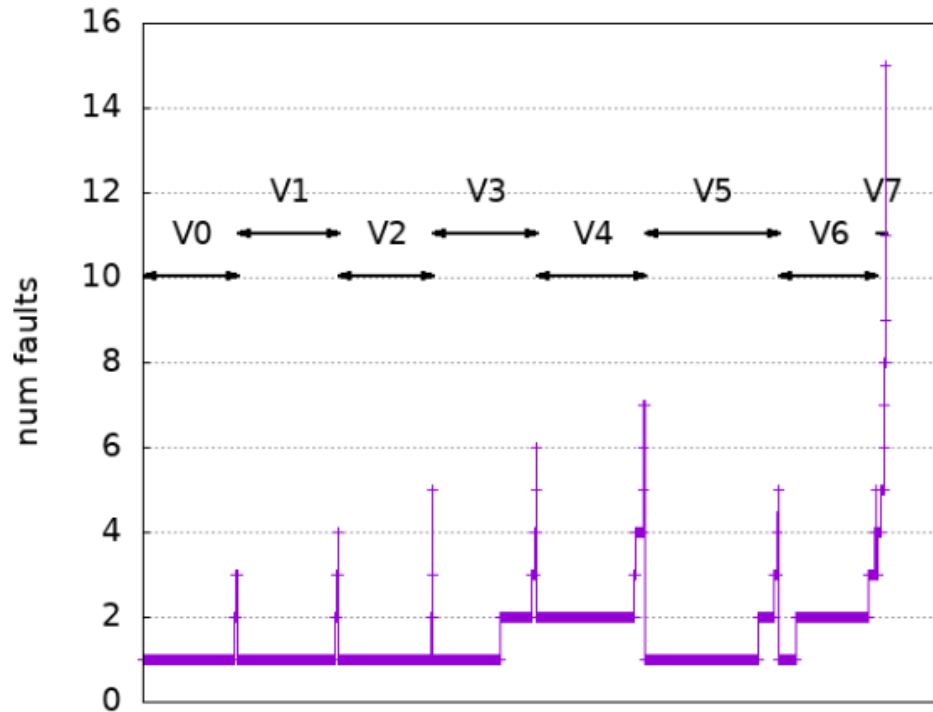
Analysis time



- Time taken to find **up to 30 attack paths**
- Xeon Gold 5218 2.3 GHz - 32 physical cores

→ **Most of the results are under the threshold of half a second**

Number of faults needed for each successful attack



- Only **1 fault**: V0 to V3, V5 and V6
- At least **2 faults**: V4
- At least **3 faults**: V7

→ **SAMVA is able to find attack paths with multiple faults when it is required**

Conclusion

- **SAMVA**

Framework based only on static analysis for determining attack paths in presence of multiple instruction-skip faults

- **Evaluation**

Attack paths found for all the 1 + 7 hardened versions of PIN code verification

- **Future work**

- Extension of supported fault models
→ “instruction replay”
- Make the link with fault injection platform

- **More details in the COSADE 23 paper !**

